

# Quantum Error Correction / CO639

lecture: 2004-02-24

Prepared by Annika Niehage

Edited by Daniel Gottesman

Guest Lecturer: Martin Rötteler

April 28, 2004

## 1 Codes over fields of arbitrary prime order

### Plan:

- higher-dimensional quantum codes
- generalize error basis (*Pauli matrices*), *Clifford group*

### Motivation:

- sometimes binary is too restrictive,  
e.g.  $[[3, 1, 2]]_2$  does not exist  
for alphabet size 3 this exists  
e.g.  $[[9, 5, 3]]_2$  does not exist  
for alphabet size 3 this guy exists
- beautiful constructions for higher-dimensional codes (e.g. *Reed Solomon codes*)

### Literature:

- [1] D. Gottesman: “*Fault-tolerant quantum computation with higher-dimensional systems,*” QCCQ ’98, quant-ph/9802007
- [2] A. Ashikhmin, M. Knill: “*Nonbinary quantum stabilizer codes,*” quant-ph/0005008

[3] M. Grassl, M. Rötteler, Th. Beth: “Efficient quantum circuits for non-qubit QECC,” quant-ph/0211014

**Example:**  $[[3, 1, 2]]_2$  does not exist ( $\rightarrow$  assignment 3)

Let’s construct a  $[[3, 1, 2]]_d$  for **any** odd  $d \in \mathbb{N}$ . Let

$$G := \left( \begin{array}{ccc|ccc} 1 & 1 & 1 & 1 & 1 & 0 \\ 2 & 0 & 0 & 0 & 1 & 1 \end{array} \right) \in \mathbb{Z}^{2 \times 6}$$

and let  $C$  be the stabilizer code given by  $G$ .

Define the **symplectic inner product** as

$$(x, y) * (a, b) = ya - xb$$

Then

$$\left( \begin{array}{ccc|ccc} 1 & 1 & 1 & 1 & 1 & 0 \end{array} \right) * \left( \begin{array}{ccc|ccc} 2 & 0 & 0 & 0 & 1 & 1 \end{array} \right) = 2 - 2 = 0$$

Show that this detects any error  $E$  of wt 1.

- One way would be to find an element  $M \in S$  s.t.  $[M, E] \neq 0$ .
- **Here:** We can assume that

$$E = \begin{pmatrix} u & 0 & 0 & | & v & 0 & 0 \\ 0 & u & 0 & | & 0 & v & 0 \\ 0 & 0 & u & | & 0 & 0 & v \end{pmatrix}, u, v \in \mathbb{Z}, (u, v) \neq (0, 0)$$

$\Rightarrow$  Compute

$$\begin{aligned} A_1 &:= G * \left( \begin{array}{ccc|ccc} u & 0 & 0 & | & v & 0 & 0 \end{array} \right) = \begin{pmatrix} u-v \\ -2v \end{pmatrix} = \underbrace{\begin{pmatrix} 1 & 1 \\ 2 & 0 \end{pmatrix}}_{\det=-2} \cdot \begin{pmatrix} -v \\ u \end{pmatrix} \\ A_2 &:= G * \left( \begin{array}{ccc|ccc} 0 & u & 0 & | & 0 & v & 0 \end{array} \right) = \begin{pmatrix} u-v \\ u \end{pmatrix} = \underbrace{\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}}_{\det=1} \cdot \begin{pmatrix} -v \\ u \end{pmatrix} \\ A_3 &:= G * \left( \begin{array}{ccc|ccc} 0 & 0 & u & | & 0 & 0 & v \end{array} \right) = \begin{pmatrix} -v \\ u \end{pmatrix} = \underbrace{\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}}_{\det=1} \cdot \begin{pmatrix} -v \\ u \end{pmatrix} \end{aligned}$$

$\Rightarrow$  can detect all *wt* 1 errors

- $\Leftrightarrow A_1, A_2, A_3$  are invertible if we compute *mod*  $d$
- $\Leftrightarrow \det A_i$  is invertible
- $\Leftrightarrow 2$  is invertible in  $\mathbb{Z}/p\mathbb{Z}$
- $\Leftrightarrow d$  is odd

**Result:** for any odd  $d \in \mathbb{N}$  there exists a  $[[3, 1, 2]]_d$  stabilizer code.

**Recall:**  $GF(p)$  is the set  $\{0, \dots, p-1\}$  equipped with usual “+”, “ $\cdot$ ” *mod*  $p$ ,  $p$  prime.

**Definition:** (**Error bases**)

Let  $p$  be a prime. Then define

$$X_\alpha := \sum_{x=0}^{p-1} |x + \alpha\rangle\langle x| \text{ for } \alpha \in GF(p)$$

$$Z_\beta := \sum_{z=0}^{p-1} \omega_p^{\beta \cdot z} |z\rangle\langle z| \text{ for } \beta \in GF(p)$$

where  $\omega_p = \exp\left(\frac{2\pi i}{p}\right)$

These operators are also called “**Weyl operators**”.

(for  $GF(q)$ ,  $q = p^r$ :

$$X_\alpha := \sum_{x \in GF(q)} |x + \alpha\rangle\langle x| \text{ for } \alpha \in GF(q)$$

$$Z_\beta := \sum_{z \in GF(q)} \omega_p^{\text{tr}(\beta \cdot z)} |z\rangle\langle z| \text{ for } \beta \in GF(p)$$

)

**Theorem:** Let  $p$  be a prime.

- (i)  $\forall \alpha, \beta : X_\alpha \cdot Z_\beta = \omega_p^{-\alpha\beta} Z_\beta \cdot X_\alpha$
- (ii)  $\forall \alpha, \alpha', \beta, \beta' : (X_\alpha \cdot Z_\beta) \cdot (X_{\alpha'} \cdot Z_{\beta'}) = \omega_p^{\alpha'\beta - \alpha\beta'} (X_{\alpha'} \cdot Z_{\beta'}) \cdot (X_\alpha \cdot Z_\beta)$
- (iii) The set  $\{X_\alpha Z_\beta : \alpha, \beta \in GF(p)\}$  is an orthonormal basis for  $\mathbb{C}^{p \times p}$  with respect to the inner product

$$\langle A, B \rangle := \frac{1}{p} \text{tr}(A^\dagger B)$$

**Proof:**

(i) We have to compute

$$\begin{aligned}
X_\alpha Z_\beta X_\alpha^{-1} &= \left( \sum_x |x + \alpha\rangle \langle x| \right) \left( \sum_z \omega_p^{\beta \cdot z} |z\rangle \langle z| \right) \left( \sum_{x'} |x'\rangle \langle x' + \alpha| \right) \\
&= \sum_x \omega_p^{\beta \cdot x} |x + \alpha\rangle \langle x + \alpha| \\
&= \sum_x \omega_p^{\beta \cdot (x - \alpha)} |x\rangle \langle x| \\
&= \sum_x \omega_p^{-\alpha \beta} \omega_p^{\beta \cdot x} |x\rangle \langle x| \\
&= \omega_p^{-\alpha \beta} \sum_x \omega_p^{\beta \cdot x} |x\rangle \langle x| \\
&= \omega_p^{-\alpha \beta} Z_\beta
\end{aligned}$$

(ii) direct consequence of (i)

(iii) To show:

(a) if  $(\alpha, \beta) \neq (\alpha', \beta')$ , then  $\text{tr}((X_\alpha Z_\beta)^\dagger (X_{\alpha'} Z_{\beta'})) \neq 0$ .

(b) This will follow if we can show that  $\text{tr}(X_\alpha Z_\beta) \neq 0$  iff  $(\alpha, \beta) \neq (0, 0)$

We compute

$$\begin{aligned}
\frac{1}{p} \text{tr}(X_\alpha Z_\beta) &= \frac{1}{p} \text{tr} \left[ \left( \sum_x |x + \alpha\rangle \langle x| \right) \left( \sum_z \omega_p^{\beta \cdot z} |z\rangle \langle z| \right) \right] \\
&= \frac{1}{p} \text{tr} \left( \sum_x \omega_p^{\beta \cdot x} |x + \alpha\rangle \langle x| \right) \\
&= \frac{1}{p} \sum_{x=0}^{p-1} \omega_p^{\beta \cdot x} \underbrace{\langle x + \alpha | x \rangle}_{\delta_{\alpha, 0}} \\
&= \delta_{\alpha, 0} \delta_{\beta, 0}
\end{aligned}$$

□

**Definition:** Let  $p$  be a prime,  $X_\alpha, Z_\beta$  as above.

For  $\vec{\alpha}, \vec{\beta} \in GF(p)^n$  define

$$\begin{aligned}
X_{\vec{\alpha}} &:= X_{\alpha_1} \otimes \cdots \otimes X_{\alpha_n} \\
Z_{\vec{\beta}} &:= Z_{\beta_1} \otimes \cdots \otimes Z_{\beta_n}
\end{aligned}$$

**Remark:** “Everything holds for the  $X_{\vec{\alpha}}, Z_{\vec{\beta}}$  as well.”

This means

- (i)  $(X_{\vec{\alpha}} \cdot Z_{\vec{\beta}}) \cdot (X_{\vec{\alpha}'} \cdot Z_{\vec{\beta}'}) = \omega_p^{\sum_{i=1}^n \alpha'_i \beta_i - \alpha_i \beta'_i} (X_{\vec{\alpha}'} \cdot Z_{\vec{\beta}'}) \cdot (X_{\vec{\alpha}} \cdot Z_{\vec{\beta}})$
- (ii) similarly
- (iii) The set  $\{X_{\vec{\alpha}} Z_{\vec{\beta}} : \vec{\alpha}, \vec{\beta} \in GF(p)^n\}$  is an ONB for  $\mathbb{C}^{p^n \times p^n}$ .

**Definition:** (Pauli group, Clifford group)

$$\begin{aligned} \mathcal{P}_{n,p} &:= \langle X_{\vec{\alpha}} Z_{\vec{\beta}} : \vec{\alpha}, \vec{\beta} \in GF(p)^n \rangle \\ \mathcal{C}_{n,p} &:= N_{U(p^n)}(\mathcal{P}_{n,p}) \end{aligned}$$

**Definition/Theorem:** (famous elements of  $\mathcal{C}_{n,p}$ )

- (i)  $DFT := \frac{1}{\sqrt{p}} \sum_{x,z} \omega_p^{xz} |z\rangle \langle x|$
- (ii)  $P := \sum_z \omega_p^{\frac{z(z-1)}{2}} |z\rangle \langle z|$
- (iii)  $ADD^{(1,2)} := \sum_{x,y} |x\rangle_1 |x+y\rangle_2 \langle x|_1 \langle y|_2$
- (iv)  $M_\gamma := \sum_{y=0}^{p-1} |\gamma y\rangle \langle y|$  for  $\gamma \in GF(p)$

**Proof:**

- (i) We have to show that  $DFT^\dagger E DFT$  is again a Pauli matrix (for all  $E \in \mathcal{P}_{1,p}$ )

$$\begin{aligned} DFT^\dagger Z_\beta DFT &= \left( \frac{1}{\sqrt{p}} \sum_{i,j} \omega_p^{-ij} |i\rangle \langle j| \right) \left( \sum_z \omega_p^{\beta z} |z\rangle \langle z| \right) \left( \frac{1}{\sqrt{p}} \sum_{k,l} \omega_p^{kl} |k\rangle \langle l| \right) \\ &= \frac{1}{p} \sum_{i,l,z} \omega_p^{(-iz + \beta z + zl)} |i\rangle \langle l| \\ &= \frac{1}{p} \sum_{i,l} \left( \underbrace{\sum_z \omega_p^{(-i+\beta+l)z}}_{\substack{= 0 \text{ if } -i+\beta+l \neq 0 \\ = p \text{ if } -i+\beta+l = 0 \Leftrightarrow i=\beta+l}} \right) |i\rangle \langle l| \\ &= \sum_l |l+\beta\rangle \langle l| \\ &= X_\beta \end{aligned}$$

Similarly, we can show:  $DFT^\dagger X_\alpha DFT = Z_\alpha^{-1} \Rightarrow DFT$  acts on  $(\alpha|\beta) \in GF(p)^2$  as the matrix

$$\begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}$$

Or the other way round with conjugation  $U \mapsto UEU^\dagger$

$$\begin{array}{ccc} X & \xrightarrow{DFT} & Z \\ Z & \mapsto & X^{-1} \end{array}$$

(Exercise): Show that

(ii)  $P$  acts like  $\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$ ,  $\begin{array}{ccc} X & \mapsto & XZ \\ Z & \mapsto & Z \end{array}$

(iii)  $M_\gamma$  acts like  $\begin{pmatrix} \gamma^{-1} & 0 \\ 0 & \gamma \end{pmatrix}$

(iv)  $ADD^{(1,2)}$  acts on  $(\alpha_1|\beta_1), (\alpha_2|\beta_2)$  as the matrix

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 \\ -1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} \alpha_1 \\ \beta_1 \\ \alpha_2 \\ \beta_2 \end{pmatrix} = \begin{pmatrix} \alpha_1 \\ \beta_1 + \beta_2 \\ -\alpha_1 + \alpha_2 \\ \beta_2 \end{pmatrix}$$

□